

“Maus programadores se preocupam com código. Bons programadores se preocupam com as estruturas de dados e seus relacionamentos” (Linus Torvalds).

Merge Sort

Paulo Ricardo Lisboa de Almeida

Merge Sort

Ideia: particionar o vetor em duas partes de (quase) mesmo tamanho

Recursivamente, para cada uma das partes, repetir o processo, até que o vetor particionado seja **trivialmente** ordenável

Juntar novamente os vetores através de uma função de merge

Merge Sort - Exemplo

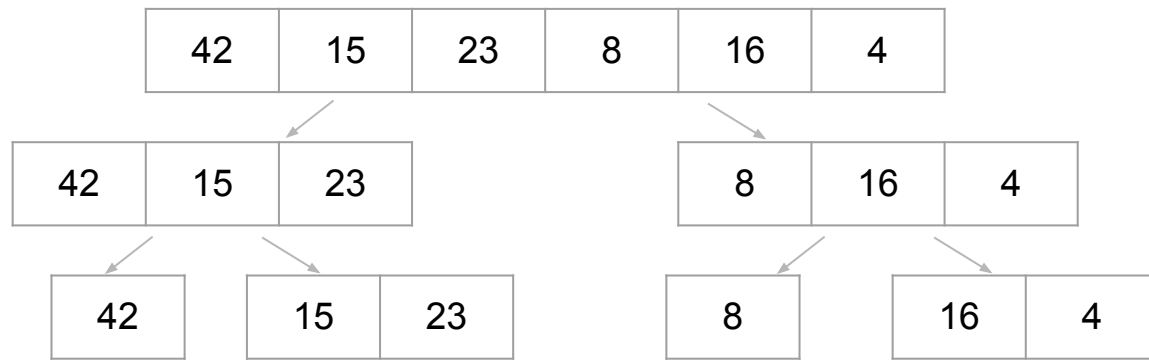
42	15	23	8	16	4
----	----	----	---	----	---

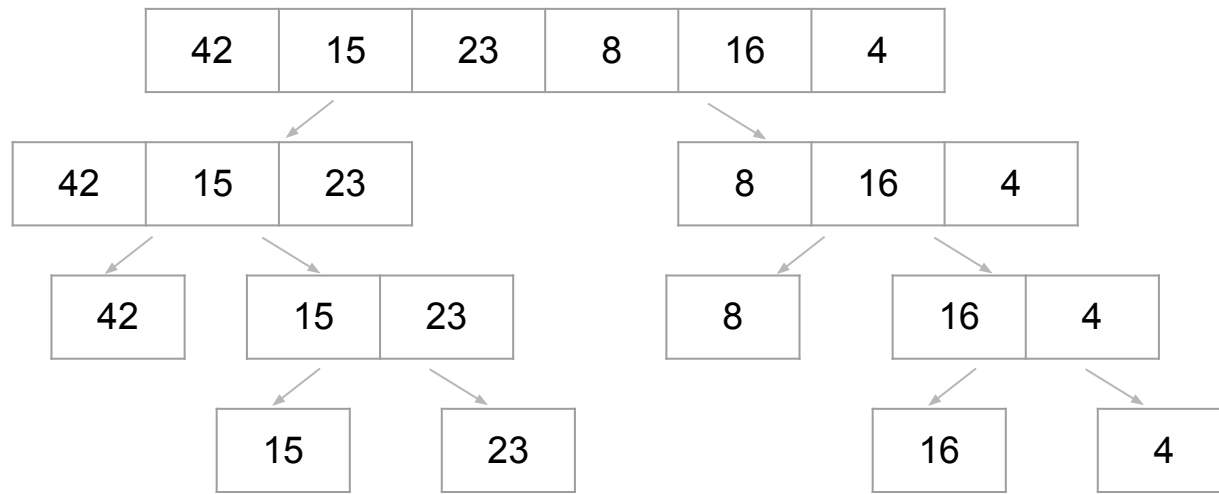
42	15	23	8	16	4
----	----	----	---	----	---

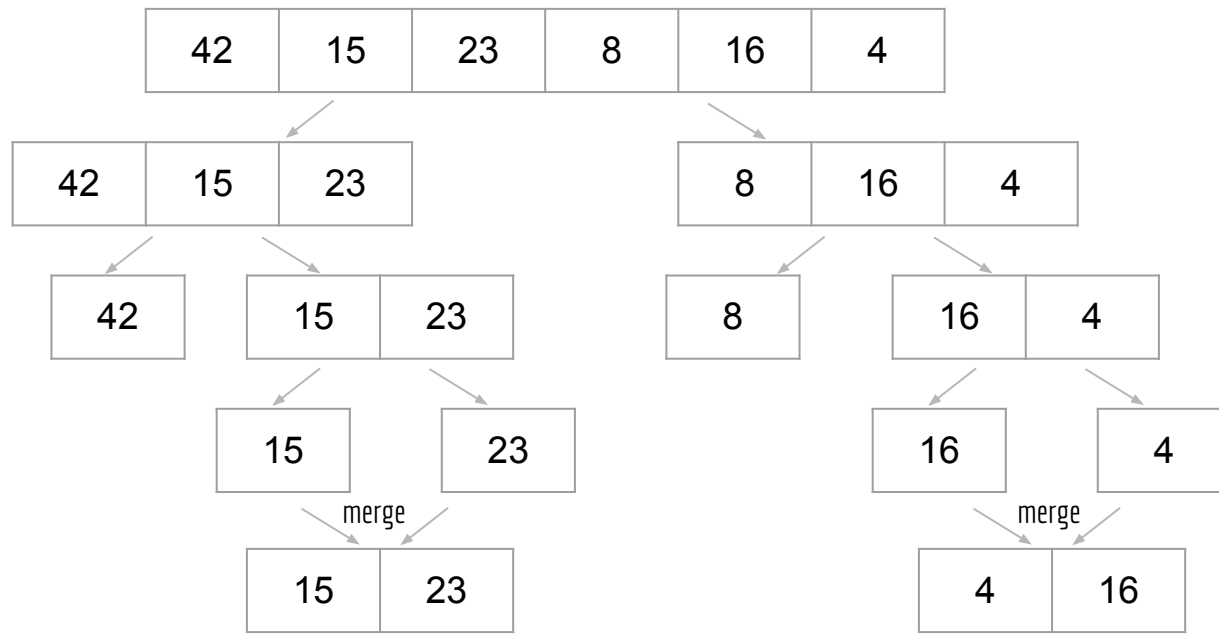


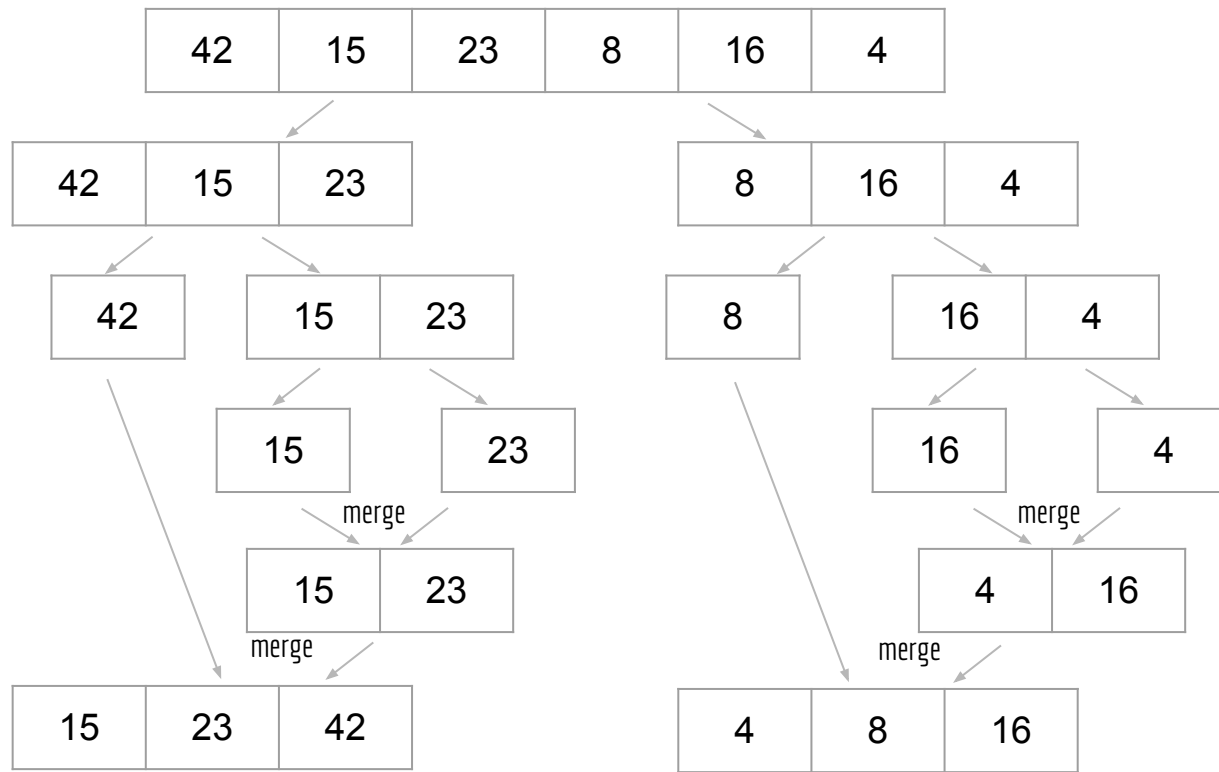
42	15	23
----	----	----

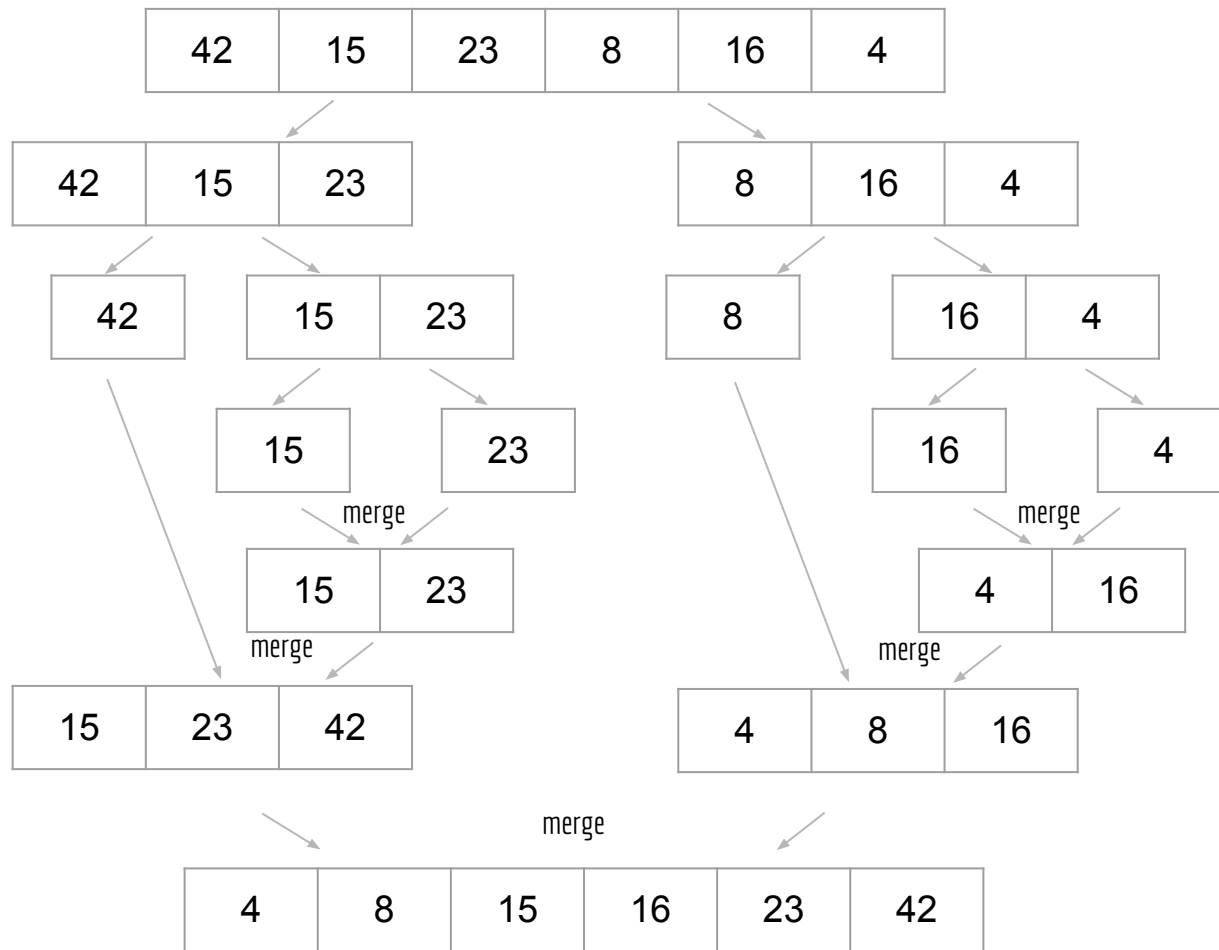
8	16	4
---	----	---











Merge-Sort

função mergeSort (v, a, b)

entrada: vetor v, indexado por [a..b]

saída: o vetor v modificado de forma que v[a..b] é um vetor ordenado.

se $a \geq b$

 retorne v
m ← $\lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Teste de Mesa

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v, a, b)

se $a \geq b$

 retorne v

m $\leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

m $\leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

m $\leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1



mergeSort		
a	b	m
1	1	

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

i	1	2	3	4
v[i]	42	15	23	8

mergeSort		
a	b	m
1	2	1

mergeSort		
a	b	m
1	1	

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1



mergeSort		
a	b	m
2	2	

i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1



mergeSort		
a	b	m
2	2	



i	1	2	3	4
v[i]	42	15	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2



mergeSort		
a	b	m
1	2	1

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	3

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	3

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	3

mergeSort		
a	b	m
3	3	

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	3

mergeSort		
a	b	m
3	3	

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	3

i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

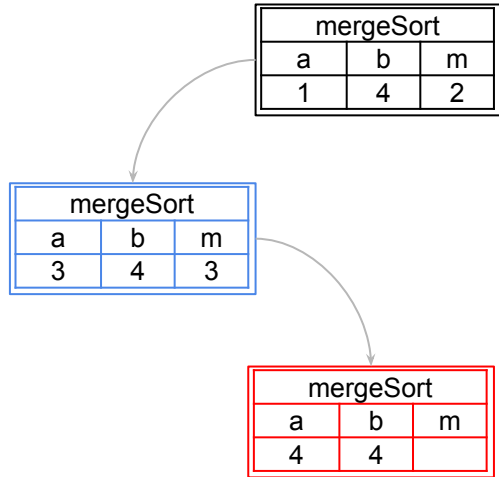
retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)



i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

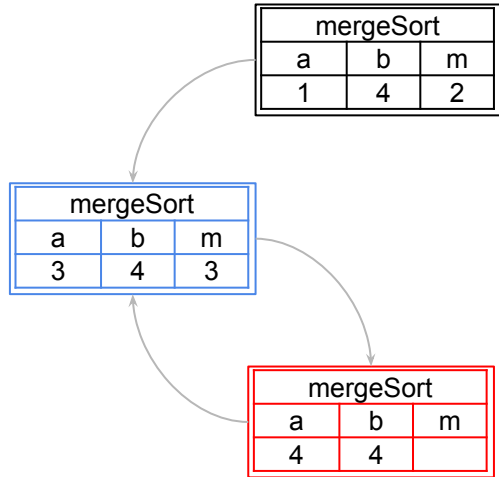
 retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)



i	1	2	3	4
v[i]	15	42	23	8

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

mergeSort		
a	b	m
3	4	3

i	1	2	3	4
v[i]	15	42	8	23

função mergeSort (v,a,b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

i	1	2	3	4
v[i]	8	15	23	42

função mergeSort (v,a,b)

se $a \geq b$

 retorne v

m $\leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

mergeSort		
a	b	m
1	4	2

i	1	2	3	4
v[i]	8	15	23	42

função mergeSort (v,a,b)

se $a \geq b$

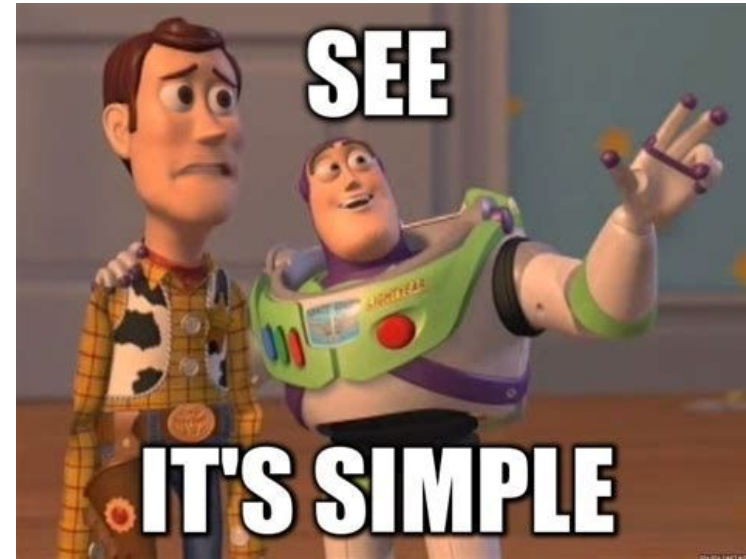
retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)



Análise

Considerando o número de comparações com elementos do vetor

$$C(n) = \dots$$

função mergeSort (v, a, b)

se $a \geq b$

 retorne v

m ← $\lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

Considerando o número de comparações com elementos do vetor

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ C(\lfloor \frac{n}{2} \rfloor) + C(\lceil \frac{n}{2} \rceil) + C_m(n), & \text{se } n > 1 \end{cases}$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

Da aula passada, o custo do merge para um vetor de tamanho n é n

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ C(\lfloor \frac{n}{2} \rfloor) + C(\lceil \frac{n}{2} \rceil) + n, & \text{se } n > 1 \end{cases}$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

Para facilitar (muito) a análise, considerando que $\lfloor \frac{n}{2} \rfloor \approx \frac{n}{2}$ e $\lceil \frac{n}{2} \rceil \approx \frac{n}{2}$

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ C(\frac{n}{2}) + C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ 2C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

$$C(n) = 2C(\frac{n}{2}) + n$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ 2C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

$$C(n) = 2C(\frac{n}{2}) + n$$

$$C(n) = 2(2C(\frac{n}{2^2}) + \frac{n}{2}) + n = 2^2C(\frac{n}{2^2}) + 2n$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ 2C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

$$C(n) = 2C(\frac{n}{2}) + n$$

$$C(n) = 2(2C(\frac{n}{2^2}) + \frac{n}{2}) + n = 2^2 C(\frac{n}{2^2}) + 2n$$

$$C(n) = 2^2(2C(\frac{n}{2^3}) + \frac{n}{2^2}) + 2n = 2^3 C(\frac{n}{2^3}) + 3n$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ 2C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

$$C(n) = 2C(\frac{n}{2}) + n$$

$$C(n) = 2(2C(\frac{n}{2^2}) + \frac{n}{2}) + n = 2^2C(\frac{n}{2^2}) + 2n$$

$$C(n) = 2^2(2C(\frac{n}{2^3}) + \frac{n}{2^2}) + 2n = 2^3C(\frac{n}{2^3}) + 3n$$

$$C(n) = 2^\mu C(\frac{n}{2^\mu}) + \mu n$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ 2C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

$$C(n) = 2^\mu C(\frac{n}{2^\mu}) + \mu n \quad \frac{n}{2^\mu} = 1 \iff \mu = \log_2 n$$

função mergeSort (v, a, b)

se a ≥ b

retorne v

m ← ⌊(a+b)/2⌋

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

retorne merge(v, a, m, b)

Análise

$$C(n) = \begin{cases} 0, & \text{se } n \leq 1, \\ 2C(\frac{n}{2}) + n, & \text{se } n > 1 \end{cases}$$

$$C(n) = 2^\mu C(\frac{n}{2^\mu}) + \mu n \quad \frac{n}{2^\mu} = 1 \iff \mu = \log_2 n$$

$$C(n) = 2^{\log_2 n} C(1) + \log_2(n)n = n \log_2 n$$

função mergeSort (v, a, b)

se $a \geq b$

retorne v

$m \leftarrow \lfloor (a+b)/2 \rfloor$

mergeSort(v, a, m)

mergeSort(v, m + 1, b)

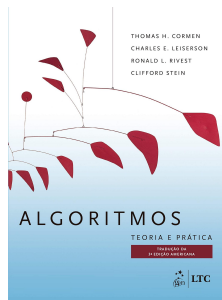
retorne merge(v, a, m, b)

Exercícios

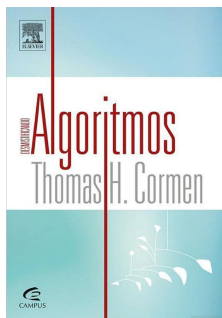
1. Plote um gráfico onde o eixo x possui valores para n , e no eixo y plote $n^2/2$ e $n \log_2 n$. Compare a diferença entre as curvas, especialmente para valores grandes de n .
2. Mostre que o custo de memória do algoritmo é aproximadamente $M(n) = n + n \cdot \log_2(n)$
3. Implemente o merge sort em C

Referências

T. Cormen, C. Leiserson,
R. Rivest, C. Stein.
Algoritmos: Teoria e
Prática. 3a ed. 2012



T. Cormen.
Desmistificando
algoritmos. 2017.

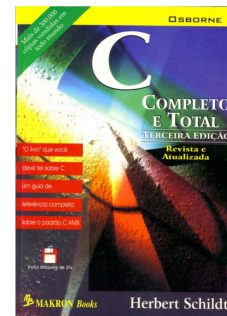


Renato Carmo. Algoritmos e
Estruturas de Dados.
www.inf.ufpr.br/renato

R. Sedgwick, K. Wayne.
Algorithms Part I. 4a ed.
2014



H. Schildt. C completo e
total. 1996



Licença

Este obra está licenciado com uma Licença [Creative Commons Atribuição 4.0 Internacional](https://creativecommons.org/licenses/by/4.0/).

